Integer-Forcing: An Algebraic Approach to Interference Management

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Motivation



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MIMO Uplink Channel



Usual Assumptions:

- Each antenna carries an independent data stream $\mathbf{x}_{\ell} \in \mathbb{C}^n$ of rate R (e.g., V-BLAST setting, cellular uplink). $\mathbf{X} = [\mathbf{x}_1 \cdots \mathbf{x}_M]^{\mathsf{T}}$.
- Usual power constraint: $\|\mathbf{x}_{\ell}\|^2 \leq n$ SNR.
- Channel model: $\mathbf{Y} = \mathbf{H}\mathbf{X} + \mathbf{Z}$
- Z is elementwise i.i.d. $\mathcal{CN}(0,1)$.
- CSIR: Only the receiver knows channel realization $\mathbf{H} \in \mathbb{C}^{M imes M}$.

- Throughout the talk, we will assume that **H** is elementwise i.i.d. Rayleigh, remains fixed throughout the block, and is only known at the receiver.
- Say that we have a scheme that achieves rate R_{scheme}(H) under channel realization H. For a target rate R, the outage probability is

$$p_{\mathsf{out}}(R) = \mathbb{P}\big(R_{\mathsf{scheme}}(\mathbf{H}) < R\big)$$

and the outage rate is

$$R_{\mathsf{out}}(\rho) = \sup \big\{ R : p_{\mathsf{out}}(R) \le \rho \big\}.$$

MIMO Uplink Channel: Joint ML Decoding



Joint Maximum Likelihood Decoding:

$$R_{\mathsf{joint}}(\mathbf{H}) = \min_{\mathcal{S} \subseteq \{1, \dots, M\}} \frac{1}{|\mathcal{S}|} \log \det \left(\mathbf{I} + \mathsf{SNR} \ \mathbf{H}_{\mathcal{S}} \mathbf{H}_{\mathcal{S}}^* \right)$$

- Corresponds to the (symmetric) outage capacity.
- Naive implementation has prohibitively high complexity.
- Of course, there are many clever ways to reduce the complexity!

MIMO Uplink Channel: Zero-Forcing and Linear MMSE



Zero-Forcing and Linear MMSE Receivers:

- Project the received signal, $\tilde{\mathbf{Y}} = \mathbf{B}\mathbf{Y}$ to eliminate interference between data streams.
- After projection, single-user decoders attempt to recover the individual data streams.
- Optimal **B** is the MMSE projection.

MIMO Uplink Channel: Zero-Forcing and Linear MMSE



Zero-Forcing and Linear MMSE Receivers:

• The m^{th} SISO decoder tries to recover \mathbf{x}_m from $\mathbf{b}_m^\mathsf{T} \mathbf{Y}$:

$$\mathsf{SINR}_{\mathsf{LMMSE},m}(\mathbf{H}) = \max_{\mathbf{b}_m} \frac{\mathsf{SNR} \ \|\mathbf{b}_m^{\mathsf{T}}\mathbf{h}_m\|^2}{1 + \mathsf{SNR} \ \sum_{\ell \neq m} \|\mathbf{b}_m^{\mathsf{T}}\mathbf{h}_\ell\|^2}$$

• Rate per user:

$$R_{\mathsf{LMMSE}}(\mathbf{H}) = \min_{m=1,\dots,M} \log \left(1 + \mathsf{SINR}_{\mathsf{LMMSE},m}(\mathbf{H}) \right)$$

MIMO Uplink Channel: Successive Interference Cancellation



Successive Interference Cancellation Receivers:

• Decode in order π . Cancel $\mathbf{x}_{\pi(1)}, \ldots, \mathbf{x}_{\pi(m-1)}$ from $\mathbf{\tilde{y}}_m$:

$$\mathsf{SINR}_{\mathsf{SIC},\pi(m)}(\mathbf{H}) = \max_{\mathbf{b}_m} \frac{\mathsf{SNR} \ \|\mathbf{b}_m^\mathsf{T} \mathbf{h}_{\pi(m)}\|^2}{1 + \mathsf{SNR} \ \sum_{\ell=m+1}^M \|\mathbf{b}_m^\mathsf{T} \mathbf{h}_{\pi(\ell)}\|^2}$$

• Rate per user:

$$R_{\text{V-BLAST II}}(\mathbf{H}) = \max_{\pi} \min_{m=1,\dots,M} \log \left(1 + \text{SINR}_{\text{SIC},\pi(m)}(\mathbf{H}) \right)$$



What if we could decode something else?

• Zero-Forcing / LMMSE: First, eliminate interference.

Then, decode individual data streams.



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• Integer-Forcing: First, decode integer-linear combinations.



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• Integer-Forcing: First, decode integer-linear combinations. Then, eliminate interference.



What if we could decode something else?

• Zero-Forcing / LMMSE: First, eliminate interference.

Then, decode individual data streams.

- Integer-Forcing: First, decode integer-linear combinations. Then, eliminate interference.
- If the integer matrix A is full rank, we can successfully recover the individual data streams.

$$\mathbf{b}_m^\mathsf{T}\mathbf{Y} = \mathbf{b}_m^\mathsf{T}\mathbf{H}\mathbf{X} + \mathbf{b}_m^\mathsf{T}\mathbf{Z}$$

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= $\mathbf{a}_m^\mathsf{T}\mathbf{X} + (\mathbf{b}_m^\mathsf{T}\mathbf{H} - \mathbf{a}_m^\mathsf{T})\mathbf{X} + \mathbf{b}_m^\mathsf{T}\mathbf{Z}$

$$\mathbf{b}_{m}^{\mathsf{T}}\mathbf{Y} = \mathbf{b}_{m}^{\mathsf{T}}\mathbf{H}\mathbf{X} + \mathbf{b}_{m}^{\mathsf{T}}\mathbf{Z}$$
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$$= \sum_{\substack{\ell=1\\ \mathsf{Codeword}}}^{M} a_{m\ell}\mathbf{x}_{\ell}^{\mathsf{T}} + \underbrace{(\mathbf{b}_{m}^{\mathsf{T}}\mathbf{H} - \mathbf{a}_{m}^{\mathsf{T}})\mathbf{X} + \mathbf{b}_{m}^{\mathsf{T}}\mathbf{Z}}_{\mathsf{Effective Noise}}$$

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- The $a_{m\ell} \in \mathbb{Z}[j]$ are Gaussian integers and the codebook should be closed under integer-linear combinations.
- We are free to choose any full-rank integer-valued matrix A.



Integer-Forcing Linear Receivers: (Zhan-Nazer-Erez-Gastpar '12) • The m^{th} SISO decoder tries to recover $\sum_{\ell} a_{m\ell} \mathbf{x}_{\ell}$ from $\mathbf{b}_m^{\mathsf{T}} \mathbf{Y}$:

$$\mathsf{SINR}_{\mathsf{IF},m}(\mathbf{H}, \mathbf{A}) = \max_{\mathbf{b}_m} \frac{\mathsf{SNR}}{\|\mathbf{b}_m\|^2 + \mathsf{SNR}\|\mathbf{b}_m^\mathsf{T}\mathbf{H} - \mathbf{a}_m^\mathsf{T}\|^2}$$

• Rate per user:

$$R_{\mathsf{IF}}(\mathbf{H}) = \max_{\mathbf{A}} \min_{m=1,\dots,M} \log^+ \left(\mathsf{SINR}_{\mathsf{IF},m}(\mathbf{H}, \mathbf{A})\right)$$

• Includes linear MMSE as a special case by setting $\mathbf{A} = \mathbf{I}$.



2 users, 2 receive antennas, Rayleigh fading, 1% outage.

Comparison: Outage Rates



4 users, 4 receive antennas, Rayleigh fading, 1% outage.

- How can we efficiently select a good integer matrix A?
- How does the performance scale with the number of users?
- How sensitive is the performance to imperfect CSIR?
- What types of SISO encoders and decoders can we use?
- What about the downlink?
- Can we move beyond this idealized problem setting?

Finding a Good Integer Matrix

$$\mathsf{SINR}_{\mathsf{IF},m}(\mathbf{H},\mathbf{A}) = \max_{\mathbf{b}_m} \frac{\mathsf{SNR}}{\|\mathbf{b}_m\|^2 + \mathsf{SNR}\|\mathbf{H}^{\mathsf{T}}\mathbf{b}_m - \mathbf{a}_m\|^2}$$

- Optimal \mathbf{b}_m is the MMSE projection.
- Plugging in and applying the Matrix Inversion Lemma, we get that

$$\mathsf{SINR}_{\mathsf{IF},m}(\mathbf{H}, \mathbf{A}) = rac{1}{\left\| \left(\mathbf{I} + \mathsf{SNR} \ \mathbf{H}^* \mathbf{H}
ight)^{-1/2} \mathbf{a}_m
ight\|^2}$$

- Finding the optimal A corresponds to finding a good lattice basis.
- This is a hard problem in general but good approximation algorithms are known, such as the LLL algorithm.
- We are currently using a slight twist: We run LLL to get a lattice basis. Then, we turn to the dual lattice and run LLL again, initializing with the first basis.



Rayleigh fading, 1% outage.

How does the performance scale with the number of users?



20 dB symmetric rate case

Rayleigh fading, 1% outage.

What is the impact of imperfect CSIR?



- Receiver only sees $\mathbf{H} + \mathbf{E}$ where \mathbf{E} is elementwise i.i.d. $\mathcal{CN}(0, \sigma^2)$.
- May result in selecting both a suboptimal integer matrix A and a suboptimal projection matrix B.

What is the impact of imperfect CSIR?



4 users, 20dB, Rayleigh fading, 1% outage.

What kinds of SISO coding schemes can be used?



- Underlying integer-forcing is the compute-and-forward framework, which is used as a black box to recover linear combinations of the messages over some finite field F_p.
- Messages are vectors over a prime-sized finite field, $\mathbf{w}_{\ell} \in \mathbb{F}_p^k$.



• Architecture is completely digital after SISO decoders.

What kinds of SISO coding schemes can be used?



- Nazer-Gastpar '11: Compute-and-forward achievability proofs via nested lattice codes.
- High-dimensional nested lattice codes lead to nice log(SINR) expressions but have high implementation complexity.
- Remember, all we actually need is that the codebook is closed under integer-linear combinations.



- What about QAM combined with a binary linear code?
- Issue: Real addition does not map well to addition over \mathbb{F}_{2^M} .

$$[x_1 + x_2] \mod 2^M \neq x_1 \oplus x_2$$

What kinds of SISO coding schemes can be used?



- What about p-ary QAM where p is prime combined with a linear code over F_p?
- Real addition maps well to addition over \mathbb{F}_p .

 $[x_1 + x_2] \bmod p = x_1 \oplus x_2$



Uncoded Integer-Forcing:

- Project by \mathbf{b}_m , take mod p, apply slicer.
- Correct if we recover $[a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mM}x_M] \mod p$ for all m.



Uncoded Integer-Forcing:

- Project by \mathbf{b}_m , take mod p, apply slicer.
- Correct if we recover $[a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mM}x_M] \mod p$ for all m.
- Is this lattice-aided reduction?



Uncoded Integer-Forcing:

- Project by \mathbf{b}_m , take mod p, apply slicer.
- Correct if we recover $[a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mM}x_M] \mod p$ for all m.
- Is this lattice-aided reduction? Nearly. We add the $\mod p$.





Coded Integer-Forcing:

- Project by \mathbf{b}_m , take mod p, apply LDPC decoding algorithm.
- Correct if we recover $[a_{m1}\mathbf{x}_1 + a_{m2}\mathbf{x}_2 + \cdots + a_{mM}\mathbf{x}_M] \mod p$ for all m.



- Lots of interesting questions on how to design low-complexity constellations and linear codes that work well for compute-and-forward.
- Several recent papers and...

- Lots of interesting questions on how to design low-complexity constellations and linear codes that work well for compute-and-forward.
- Several recent papers and...
- Krishna's talk coming up next!

MIMO Downlink Channel



Capacity region is known. Requires dirty-paper coding.
 Caire-Shamai '03, Vishwanath-Jindal-Goldsmith '04,
 Viswanath-Tse '03, Yu-Cioffi '04, Weingarten-Steinberg-Shamai '06.

MIMO Downlink Channel: Zero-Forcing



Zero-Forcing Beamforming:

• Use beamforming matrix **B** to eliminate interference between data streams.



- Use beamforming matrix **B** to create an integer-valued effective channel **A**.
- Decode linear combinations with $q_{m\ell} = [a_{m\ell}] \mod p$.



Integer-Forcing Beamforming: (Hong-Caire '12,'13)

- Use beamforming matrix **B** to create an integer-valued effective channel **A**.
- Decode linear combinations with $q_{m\ell} = [a_{m\ell}] \mod p$. Pre-invert $\mathbf{Q} = [\mathbf{A}] \mod p$ and decode messages.



Integer-Forcing Beamforming: (Hong-Caire '12,'13)

- Use beamforming matrix **B** to create an integer-valued effective channel **A**.
- Decode linear combinations with q_{mℓ} = [a_{mℓ}] mod p.
 Pre-invert Q = [A] mod p and decode messages.
- In very recent work, we have shown that uplink-downlink duality holds for integer-forcing. He-Nazer-Shamai '14

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- What about successive cancellation for integer-forcing?
- Ordentlich-Erez-Nazer '13: Framework for IF-SIC. Exact optimality if CSIT is available. Rate points tend to lie very close to the symmetric capacity.

- Low-complexity constellations and codes.
- New algorithms for finding integer matrix A.
- Synchronization.
- What if the channel realization changes over the coding blocklength? (e.g., OFDM)
- How should we include rate adaptation?
- What does this mean for user selection?
- With Behnaam, Krishna, and students, we are working towards a WARP implementation.
- Any others?